

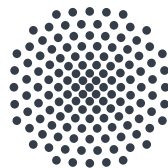
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# From Shapes to Shapes

Inferring SHACL Shapes for Results of  
SPARQL CONSTRUCT Queries

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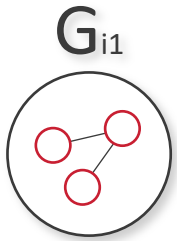
*Philipp Seifer* · *Daniel Hernández* · *Ralf Lämmel* · *Steffen Staab*



University of Stuttgart  
Germany

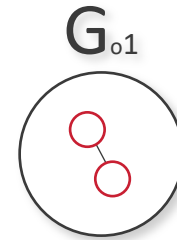
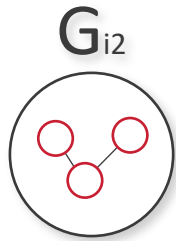
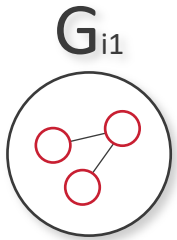


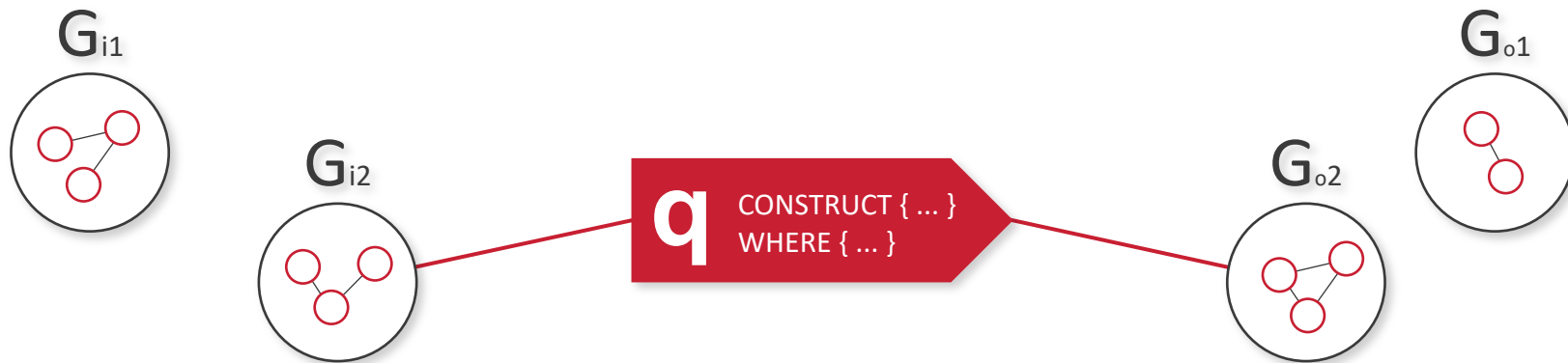
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WHERE { ... }

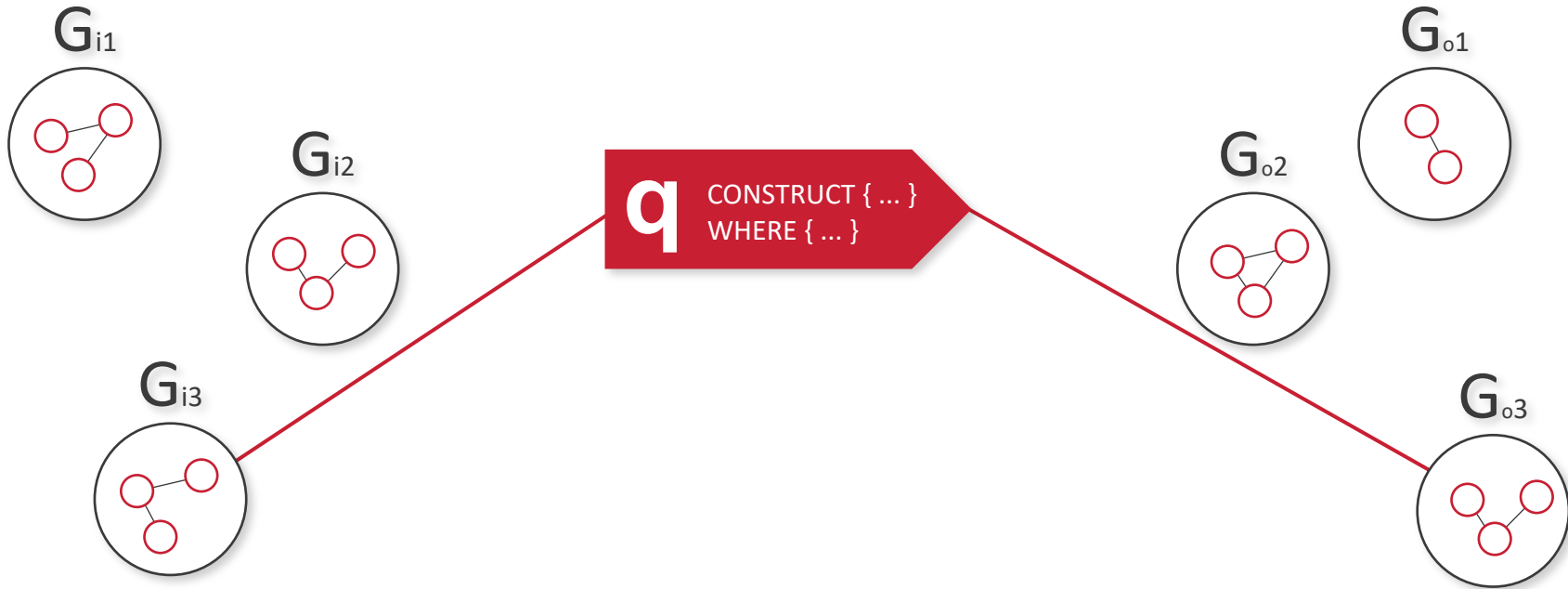


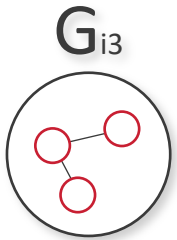
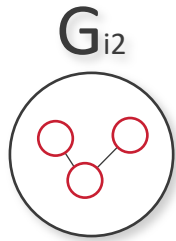
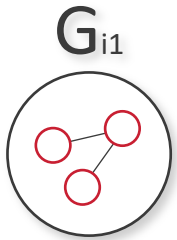
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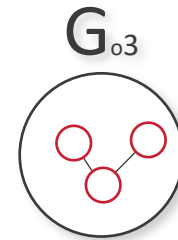
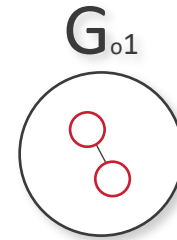






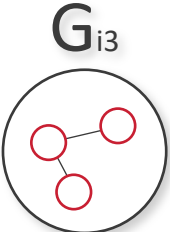
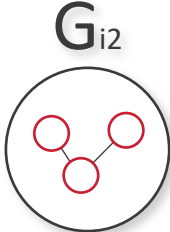
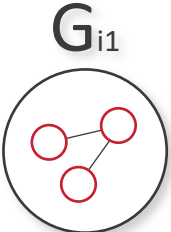


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WHERE { ... }

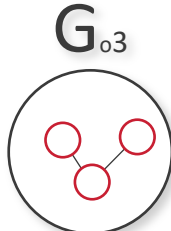
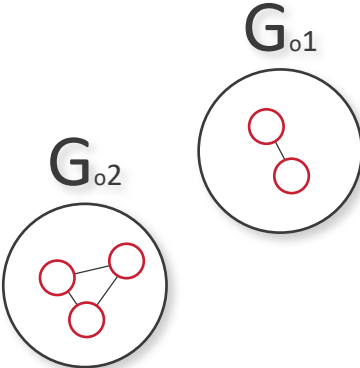




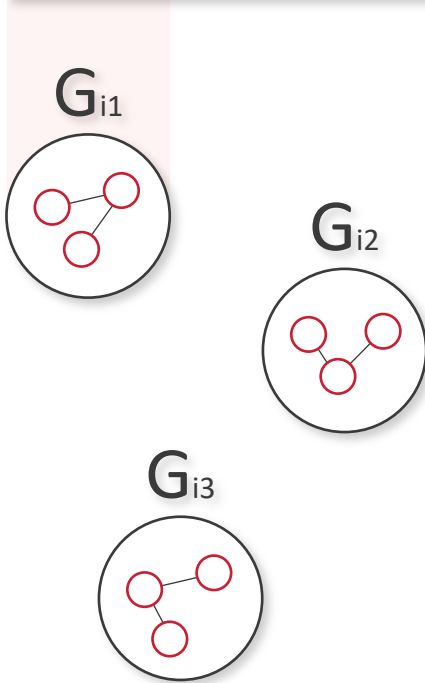
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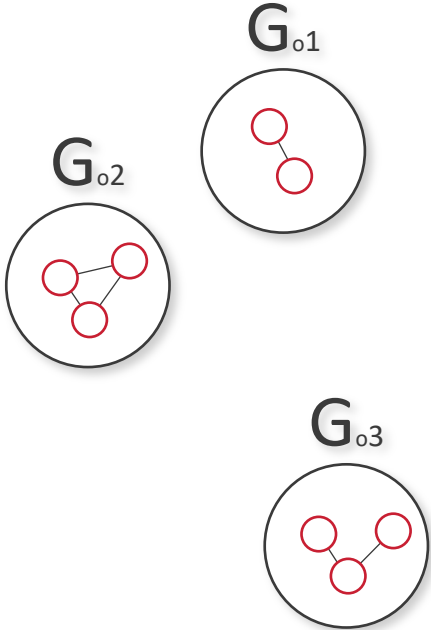
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WHERE { ... }



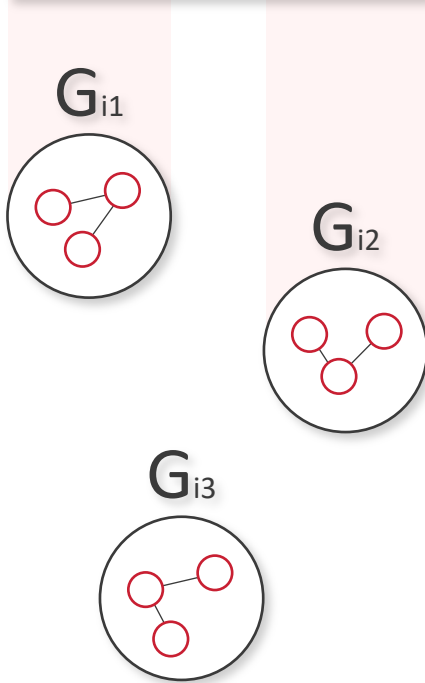
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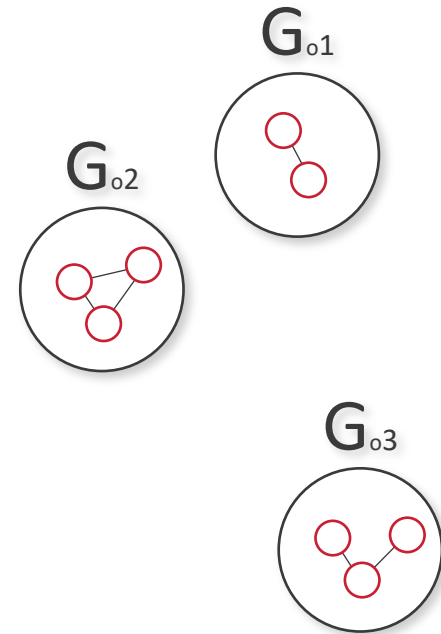
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WHERE { ... }



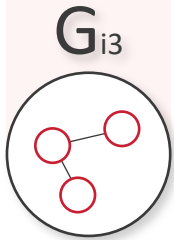
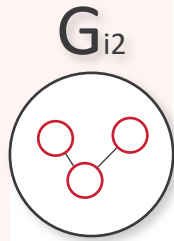
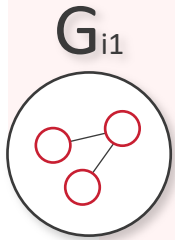
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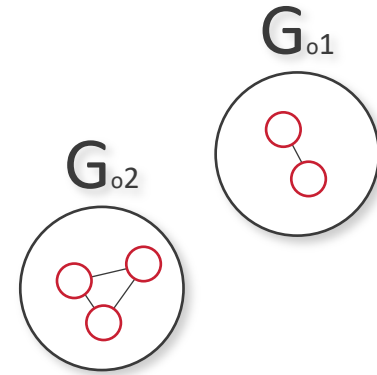
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WHERE { ... }



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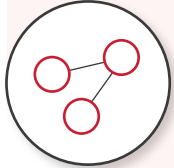


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WHERE { ... }

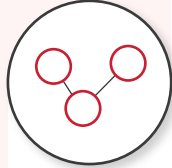


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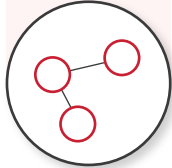
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$G_{i2}$



$G_{i3}$

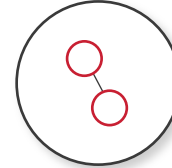


From Shapes  
to Shapes

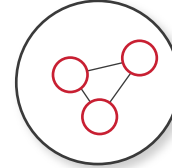
**q**

CONSTRUCT { ... }  
WHERE { ... }

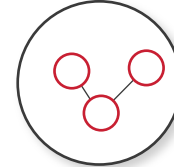
$G_{o1}$



$G_{o2}$

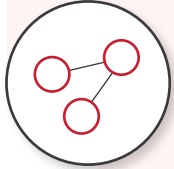


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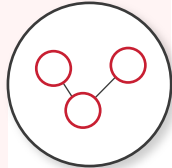


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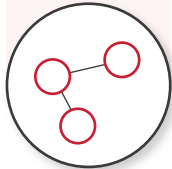
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$G_{i2}$



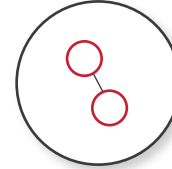
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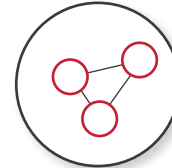
From Shapes  
to Shapes

**q** CONSTRUCT { ... }  
WHERE { ... }

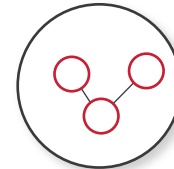
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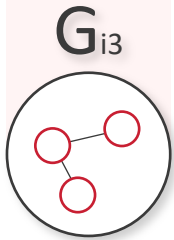
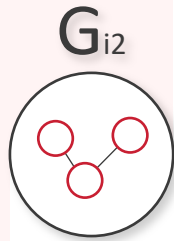
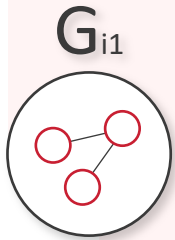
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$G_{o3}$

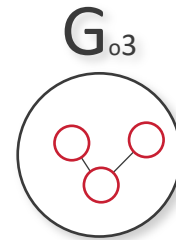
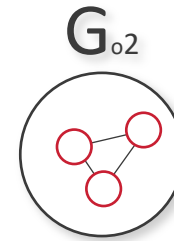
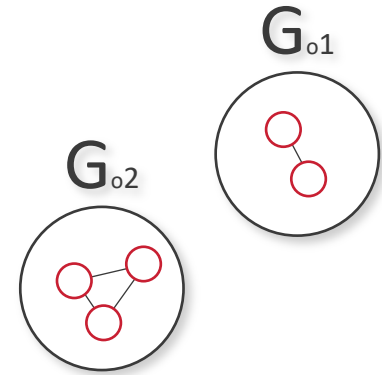


$S_{in}$

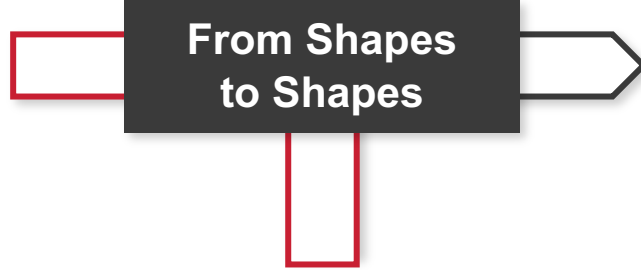
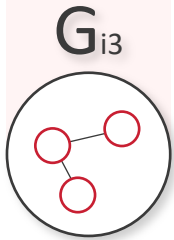
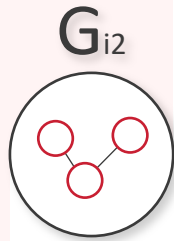
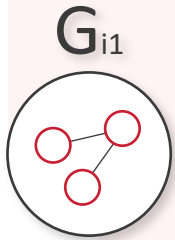


From Shapes  
to Shapes

**q** CONSTRUCT { ... }  
WHERE { ... }

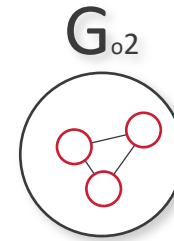
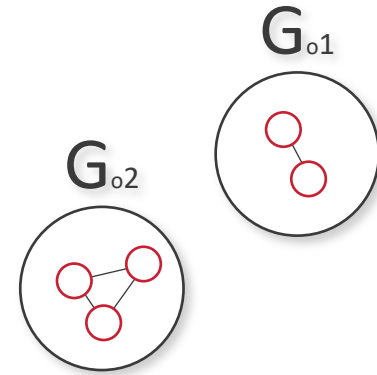


$S_{in}$



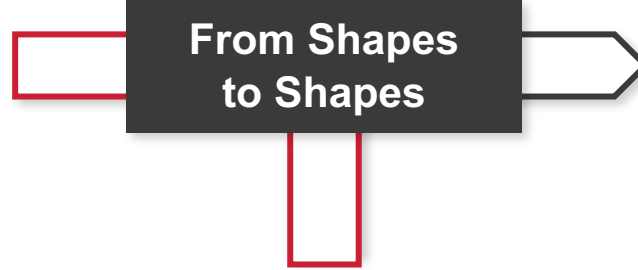
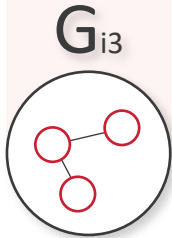
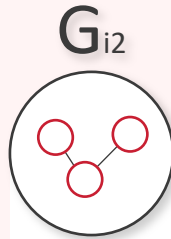
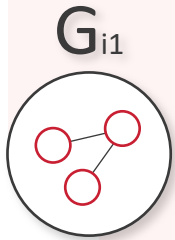
**q** CONSTRUCT { ... }  
WHERE { ... }

$S_{out}$



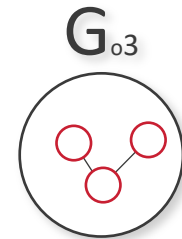
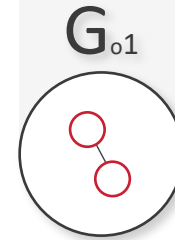


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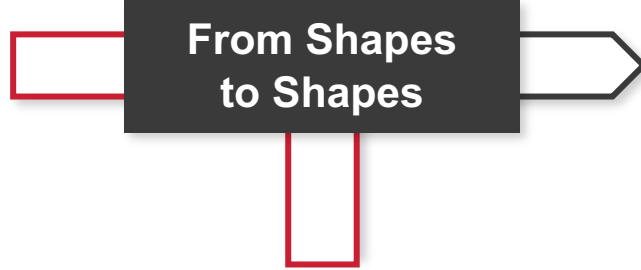
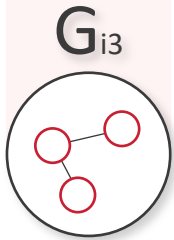
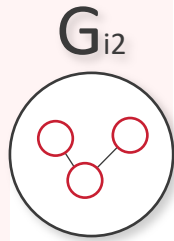
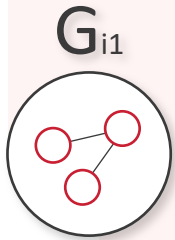


**q** CONSTRUCT { ... }  
WHERE { ... }

$S_{out}$

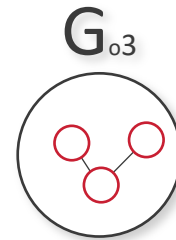
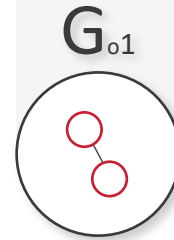
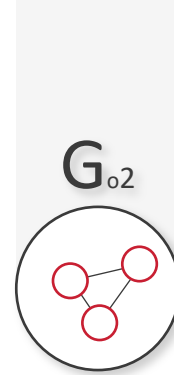


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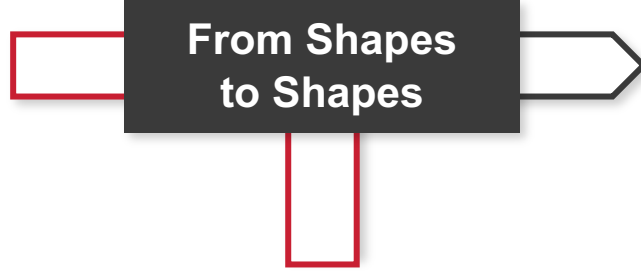
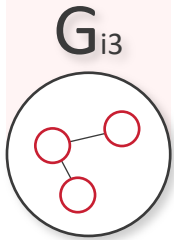
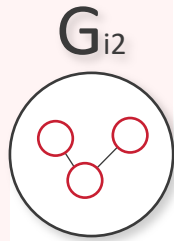
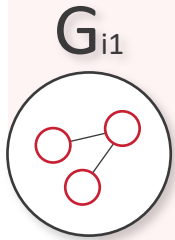


**q** CONSTRUCT { ... }  
WHERE { ... }

$S_{out}$

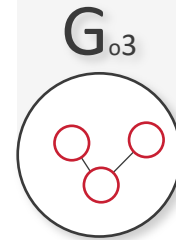
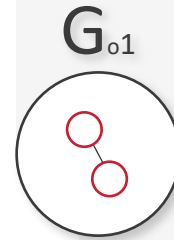
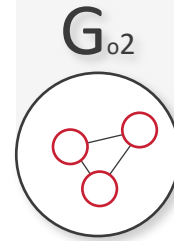


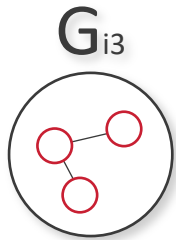
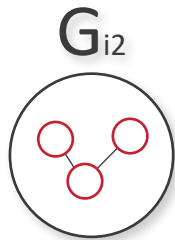
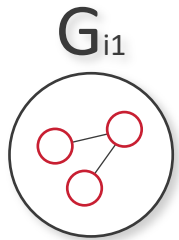
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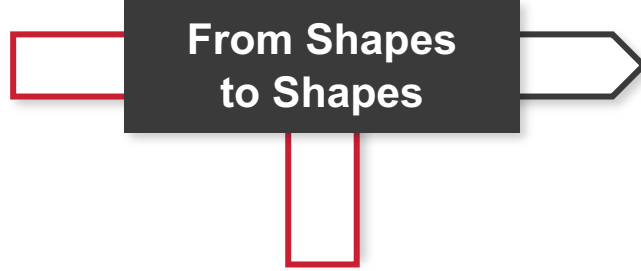
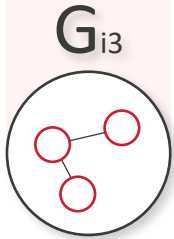
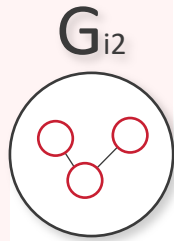
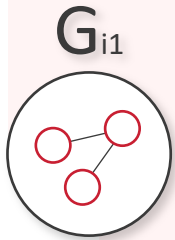
**q** CONSTRUCT { ... }  
WHERE { ... }

$S_{out}$



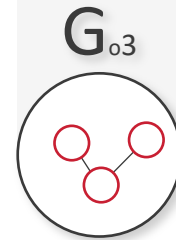
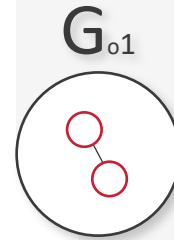
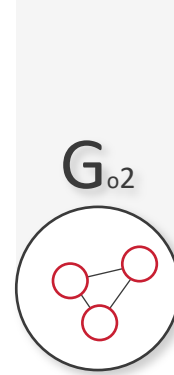


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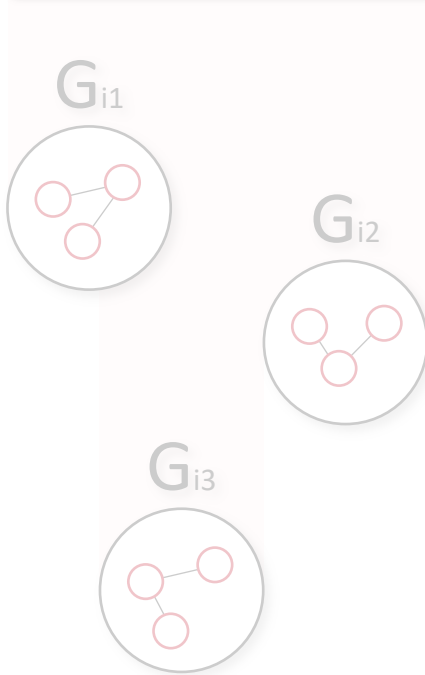


**q** CONSTRUCT { ... }  
WHERE { ... }

$S_{out}$



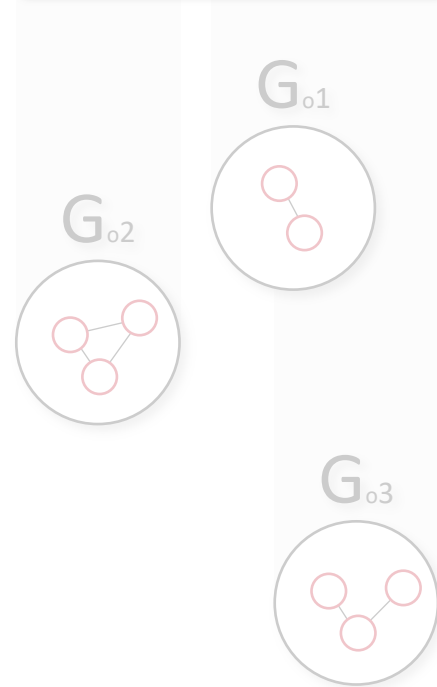
$S_{in}$



From Shapes  
to Shapes

**q** CONSTRUCT { ... }  
WHERE { ... }

$S_{out}$



## From Shapes to Shapes

```
CONSTRUCT {
```

```
} WHERE {  
  ?w :p ?y .  
  ?y a :B .
```

```
}
```

## From Shapes to Shapes

```
CONSTRUCT {
```

```
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```



## From Shapes to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{in}$

“Every B is an E.”

From Shapes  
to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{in}$

“Every B is an E.”

From Shapes  
to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{in}$

“Every B is an E.”

From Shapes  
to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{in}$

“Every B is an E.”

From Shapes  
to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{in}$

From Shapes  
to Shapes

“Every B is an E.”

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

“?y”  $\subseteq$  “?z”

$S_{in}$

From Shapes  
to Shapes

“Every B is an E.”

“?y”  $\subseteq$  “?z”

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{in}$

From Shapes  
to Shapes

“Every B is an E.”

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CONSTRUCT {  
  ?y a :F .  
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  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

“?y”  $\subseteq$  “?z”



$S_{in}$

“Every B is an E.”

“?y”  $\subseteq$  “?z”

From Shapes  
to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{out}$

“Every F is a G.”

$S_{in}$

“Every B is an E.”

From Shapes  
to Shapes

CONSTRUCT {

?y a :F .

?y :r ?z .

?z a :G

} WHERE {

?w :p ?y .

?y a :B .

?x :p ?z .

?z a :E

}

$S_{out}$

“Every F is a G.”

“Every F has at least  
one r-edge to a G.”

$S_{in}$

“Every B is an E.”

From Shapes  
to Shapes

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$S_{out}$

“Every F is a G.”

“Every F has at least  
one r-edge to a G.”

...

“Every **B** is an **E**.”

“Every B is an E.”

**Bogaerts et. al, 2022.**

*SHACL: A Description Logic in Disguise.*

LPNMR. Springer

“Every **B** is an **E**.”

**Bogaerts et. al**, 2022.

*SHACL: A Description Logic in Disguise.*

LPNMR. Springer

**B**  $\sqsubseteq$  **E**

*ALCHOI*

$\Sigma :=$

$$\Sigma := S_{in}$$



$$\Sigma := S_{in} \cup Q_{infer}$$

$$\Sigma := S_{\text{in}} \cup Q_{\text{infer}}$$

if  $\Sigma \vdash s$

$$\Sigma := S_{\text{in}} \cup Q_{\text{infer}}$$

if  $\Sigma \vdash s$  then  $s \in S_{\text{out}}$

*(modulo namespaces)*

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \sqsubseteq E$

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \subseteq E$

X encoding ?x  
Y encoding ?y  
Z encoding ?z

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \sqsubseteq E$

$X \equiv \exists p.Z$

$W \equiv \exists p.Y$

$Z \equiv E \sqcap (\exists p \neg .X)$

$Y \equiv B \sqcap (\exists p \neg .W)$

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \sqsubseteq E$

$X \equiv \exists p.Z$

$W \equiv \exists p.Y$

$Z \equiv E \cap (\exists p \neg .X)$

$Y \equiv B \cap (\exists p \neg .W)$

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \sqsubseteq E$

$X \equiv \exists p.Z$

$W \equiv \exists p.Y$

$G \equiv Z$

$\exists r.Z \equiv Y$

$\exists r.Y \equiv Z$

$Z \equiv E \cap (\exists p.X)$

$Y \equiv B \cap (\exists p.W)$

...

$F \equiv Y$

$\exists r.T \equiv Y \cap (\exists r.Z)$

$\exists r.T \equiv Z \cap (\exists p.Y)$

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```



$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$X \equiv \exists p. Z$$

$$W \equiv \exists p. Y$$

$$G \equiv Z$$

$$\exists r. Z \equiv Y$$

$$\exists r. Y \equiv Z$$

...

$$Z \equiv E \cap (\exists p. X)$$

$$Y \equiv B \cap (\exists p. W)$$

$$F \equiv Y$$

$$\exists r. T \equiv Y \cap (\exists r. Z)$$

$$\exists r. T \equiv Z \cap (\exists p. Y)$$

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

**Definition 16** (CWA-encoding). The *CWA-encoding* for a SCCQ  $q = (H \leftarrow P)$ , denoted  $\text{CWA}(q)$ , is the minimal set of  $\mathcal{ALCHOI}$  axioms including:

1. For each concept name  $A$  in  $P$ ,  $\dot{A} \equiv A \sqcap \bigsqcup_{u:A \in P} C_u$ .
2. For each concept name  $A$  in  $H$ ,  $\ddot{A} \equiv \bigsqcup_{u:A \in H} C_u$ .
3. For each variable  $x$  in  $\text{var}(q)$  the axiom

$$V_x \sqsubseteq \prod_{x:A \in P} A \sqcap \prod_{(x,u):p \in P} \exists p. C_u \sqcap \prod_{(u,x):p \in P} \exists p^-. C_u,$$

and if  $\text{vcg}(P)$  is acyclic w.r.t  $x$ , then also the axiom

$$V_x \supseteq \prod_{x:A \in P} A \sqcap \prod_{(x,u):p \in P} \exists p. C_u \sqcap \prod_{(u,x):p \in P} \exists p^-. C_u.$$

4. For each role name  $p$  in pattern  $P$  the axioms

$$\exists \dot{p}. C_v \equiv \bigsqcup_{(u,v):p \in P} C_u, \quad \exists \dot{p}. \top \equiv \bigsqcup_{(u,v):p \in P} C_u \sqcap \exists \dot{p}. C_v,$$

$$\exists \dot{p}^-. C_u \equiv \bigsqcup_{(u,v):p \in P} C_v, \quad \exists \dot{p}^-. \top \equiv \bigsqcup_{(u,v):p \in P} C_v \sqcap \exists \dot{p}^-. C_u.$$

5. For each role name  $p$  in template  $H$  the axioms

$$\exists \ddot{p}. C_v \equiv \bigsqcup_{(u,v):p \in H} C_u, \quad \exists \ddot{p}. \top \equiv \bigsqcup_{(u,v):p \in H} C_u \sqcap \exists \ddot{p}. C_v,$$

$$\exists \ddot{p}^-. C_u \equiv \bigsqcup_{(u,v):p \in H} C_v, \quad \exists \ddot{p}^-. \top \equiv \bigsqcup_{(u,v):p \in H} C_v \sqcap \exists \ddot{p}^-. C_u.$$

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \sqsubseteq E$

$X \equiv \exists p.Z$

$W \equiv \exists p.Y$

$G \equiv Z$

$\exists r.Z \equiv Y$

$\exists r.Y \equiv Z$

$Z \equiv E \cap (\exists p.X)$

$Y \equiv B \cap (\exists p.W)$

...

$F \equiv Y$

$\exists r.T \equiv Y \cap (\exists r.Z)$

$\exists r.T \equiv Z \cap (\exists p.Y)$

```
CONSTRUCT {  
  ?y a :F .  
  ?y :r ?z .  
  ?z a :G  
} WHERE {  
  ?w :p ?y .  
  ?y a :B .  
  ?x :p ?z .  
  ?z a :E  
}
```

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$B \sqsubseteq E$

$Y \sqsubseteq Z \quad ?$

$X \equiv \exists p. Z$

$W \equiv \exists p. Y$

$Z \equiv E \cap (\exists p. X)$

$Y \equiv B \cap (\exists p. W)$

...

$G \equiv Z$

$\exists r. Z \equiv Y$

$\exists r. Y \equiv Z$

$F \equiv Y$

$\exists r. T \equiv Y \cap (\exists r. Z)$

$\exists r. T \equiv Z \cap (\exists p. Y)$

```

CONSTRUCT {
  ?y a :F .
  ?y :r ?z .
  ?z a :G
} WHERE {
  ?w :p ?y .
  ?y a :B .
  ?x :p ?z .
  ?z a :E
}

```

```
CONSTRUCT {
```

```
  ...
```

```
} WHERE {
```

```
  ?w :p ?y .
```

```
  ?y a :B .
```

```
  ?x :p ?z .
```

```
  ?z a :E .
```

```
}
```

```
CONSTRUCT {
```

```
...
```

```
} WHERE {
```

```
?w :p ?y .
```

```
?y a :B .
```

```
?x :p ?z .
```

```
?z a :E .
```

```
}
```

```
?w :p ?y .
```

```
?y a :B .
```

1

```
CONSTRUCT {
```

```
...
```

```
} WHERE {
```

```
?w :p ?y .
```

```
?y a :B .
```

```
?x :p ?z .
```

```
?z a :E .
```

```
}
```

```
?w :p ?y .
```

```
?y a :B .
```

1

```
?x :p ?z .
```

```
?z a :E .
```

2



?w :p ?y .  
?y a :B .

Find a mapping **h** of query variables, such that **one component** is subset of **the other**.

?x :p ?z .  
?z a :E .

?w :p ?y .  
?y a :B .

Find a mapping **h** of query variables, such that **one component** is subset of **the other**.

?x :p ?z .  
?z a :E .

Then a mapping such as **h(x) = w** implies **W ⊆ X**.

?w :p ?y .  
?y a :B .

?x :p ?z .  
?z a :E .

?w :p ?y .  
?y a :B .

B  $\subseteq$  E

?x :p ?z .  
?z a :E .

?w :p ?y .  
?y a :B .

$B \subseteq E$

?w :p ?y .  
?y a :B .  
**?y a :E .**

?x :p ?z .  
?z a :E .

?x :p ?z .  
?z a :E .

?w :p ?y .

?y a :B .

?y a :E .

?x :p ?z .

?z a :E .

**$h(x) = w$**

**$h(z) = y$**

?w :p ?y .  
?y a :B .  
**?y a :E .**

?x :p ?z .  
?z a :E .

**h(x) = w**  
**h(z) = y**

**?w** :p **?y** .  
**?y** a :E .

?w :p ?y .  
?y a :B .  
**?y a :E .**

?w :p ?y .  
?y a :B .  
**?y a :E .**

UI

?x :p ?z .  
?z a :E .

**h(x) = w**  
**h(z) = y**

**?w** :p **?y** .  
**?y** a :E .



$$h(x) = w$$
$$h(z) = y$$

?w :p ?y .  
?y a :B .  
?y a :E .

U

$$W \subseteq X$$
$$Y \subseteq Z$$

?w :p ?y .  
?y a :E .

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$\begin{array}{l} W \sqsubseteq X \\ Y \sqsubseteq Z \end{array}$$

$$X \equiv \exists p. Z$$

$$W \equiv \exists p. Y$$

$$Z \equiv E \cap (\exists p. X)$$

$$Y \equiv B \cap (\exists p. W)$$

...

$$G \equiv Z$$

$$\exists r. Z \equiv Y$$

$$\exists r. Y \equiv Z$$

$$F \equiv Y$$

$$\exists r. T \equiv Y \cap (\exists r. Z)$$

$$\exists r. T \equiv Z \cap (\exists p. Y)$$

```

CONSTRUCT {
  ?y a :F .
  ?y :r ?z .
  ?z a :G
} WHERE {
  ?w :p ?y .
  ?y a :B .
  ?x :p ?z .
  ?z a :E
}

```

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$W \sqsubseteq X$$

$$Y \sqsubseteq Z$$

$$X \equiv \exists p. Z$$

$$Z \equiv E \cap (\exists p. X)$$

$$W \equiv \exists p. Y$$

$$Y \equiv B \cap (\exists p. W)$$

...

$$G \equiv Z$$

$$F \equiv Y$$

$$\exists r. Z \equiv Y$$

$$\exists r. T \equiv Y \cap (\exists r. Z)$$

$$\exists r. Y \equiv Z$$

$$\exists r. T \equiv Z \cap (\exists p. Y)$$

$S_{\text{out}}$



$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$W \sqsubseteq X$$

$$Y \sqsubseteq Z$$

$$X \equiv \exists p.Z$$

$$Z \equiv E \cap (\exists p \neg.X)$$

$$W \equiv \exists p.Y$$

$$Y \equiv B \cap (\exists p \neg.W)$$

...

$$G \equiv Z$$

$$F \equiv Y$$

$$\exists r.Z \equiv Y$$

$$\exists r.T \equiv Y \cap (\exists r.Z)$$

$$\exists r \neg.Y \equiv Z$$

$$\exists r \neg.T \equiv Z \cap (\exists p \neg.Y)$$

┌

$S_{\text{out}}$

$$F \sqsubseteq G$$

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$W \sqsubseteq X$$

$$Y \sqsubseteq Z$$

$$X \equiv \exists p.Z$$

$$Z \equiv E \cap (\exists p.X)$$

$$W \equiv \exists p.Y$$

$$Y \equiv B \cap (\exists p.W)$$

...

$$G \equiv Z$$

$$F \equiv Y$$

$$\exists r.Z \equiv Y$$

$$\exists r.T \equiv Y \cap (\exists r.Z)$$

$$\exists r.Y \equiv Z$$

$$\exists r.T \equiv Z \cap (\exists p.Y)$$

$$S_{\text{out}}$$

$$F \sqsubseteq G$$

$$F \sqsubseteq \exists r.G$$



$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$W \sqsubseteq X$$

$$Y \sqsubseteq Z$$

$$X \equiv \exists p. Z$$

$$Z \equiv E \cap (\exists p. X)$$

$$W \equiv \exists p. Y$$

$$Y \equiv B \cap (\exists p. W)$$

...

$$G \equiv Z$$

$$F \equiv Y$$

$$\exists r. Z \equiv Y$$

$$\exists r. T \equiv Y \cap (\exists r. Z)$$

$$\exists r. Y \equiv Z$$

$$\exists r. T \equiv Z \cap (\exists p. Y)$$

$S_{\text{out}}$

$$F \sqsubseteq G$$

$$F \sqsubseteq \exists r. G$$

$$\exists r. T \sqsubseteq F$$

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$W \sqsubseteq X$$

$$Y \sqsubseteq Z$$

$$X \equiv \exists p. Z$$

$$Z \equiv E \cap (\exists p. X)$$

$$W \equiv \exists p. Y$$

$$Y \equiv B \cap (\exists p. W)$$

...

$$G \equiv Z$$

$$F \equiv Y$$

$$\exists r. Z \equiv Y$$

$$\exists r. T \equiv Y \cap (\exists r. Z)$$

$$\exists r. Y \equiv Z$$

$$\exists r. T \equiv Z \cap (\exists p. Y)$$

$$S_{\text{out}}$$

$$F \sqsubseteq G$$

$$F \sqsubseteq \exists r. G$$

$$\exists r. T \sqsubseteq F$$

$$G \sqsubseteq \forall r. G$$

$$\Sigma := S_{\text{in}} \cup q_{\text{infer}}$$

$$B \sqsubseteq E$$

$$W \sqsubseteq X$$

$$Y \sqsubseteq Z$$

$$X \equiv \exists p. Z$$

$$Z \equiv E \cap (\exists p. X)$$

$$W \equiv \exists p. Y$$

$$Y \equiv B \cap (\exists p. W)$$

...

$$G \equiv Z$$

$$F \equiv Y$$

$$\exists r. Z \equiv Y$$

$$\exists r. T \equiv Y \cap (\exists r. Z)$$

$$\exists r. Y \equiv Z$$

$$\exists r. T \equiv Z \cap (\exists p. Y)$$

$$S_{\text{out}}$$

$$F \sqsubseteq G$$

$$F \sqsubseteq \exists r. G$$

$$\exists r. T \sqsubseteq F$$

$$G \sqsubseteq \forall r. G$$

...



Try out **Shapes 2 Shapes** at

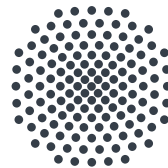
<https://github.com/softlang/s2s>



# From Shapes to Shapes

Inferring SHACL Shapes for Results of  
SPARQL CONSTRUCT Queries

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